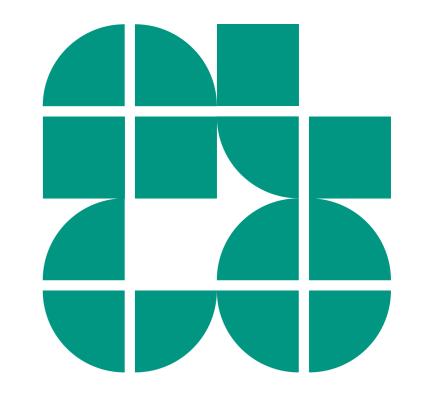


Karlsruher Institut für Technologie



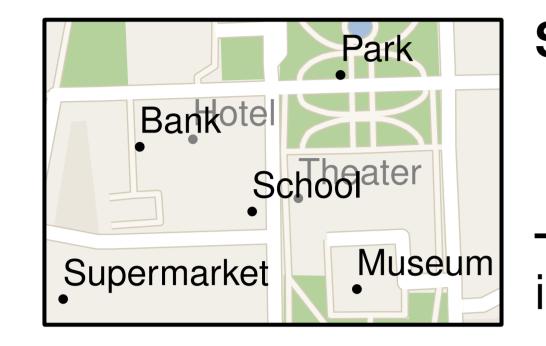
Algorithmic Cartography – Labeling

Benjamin Niedermann

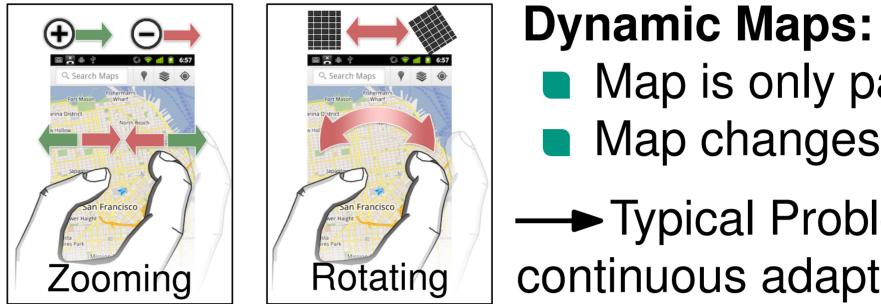
Overview

Algorithmic Cartography: With increasing popularity of interactive maps, e.g., as digital globes or on mobile devices, algorithms for creating maps become more and more important.

Labeling: Non-overlapping labels for various map features is an important cartographic problem.









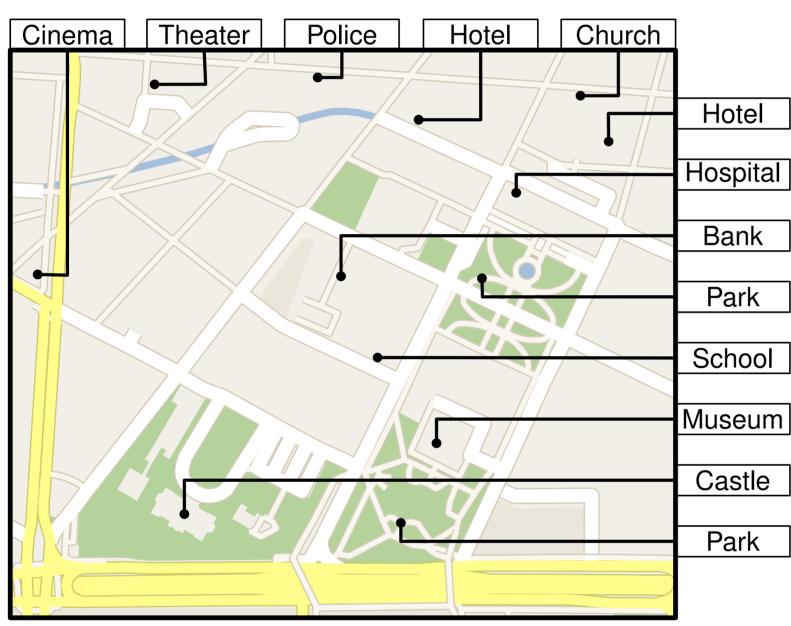
- Map does not change over time.
- independent set of labels.

- Map is only partly visible. Map changes over time.
- Typical Problem: Find consistent & continuous adaption of the labeling.

Methodology: Computational geometry, graph drawing, combinatorical optimization, algorithm engineering. **Goal:** Algorithms with guarantees for quality and performance.

Research Interests – Two Examples

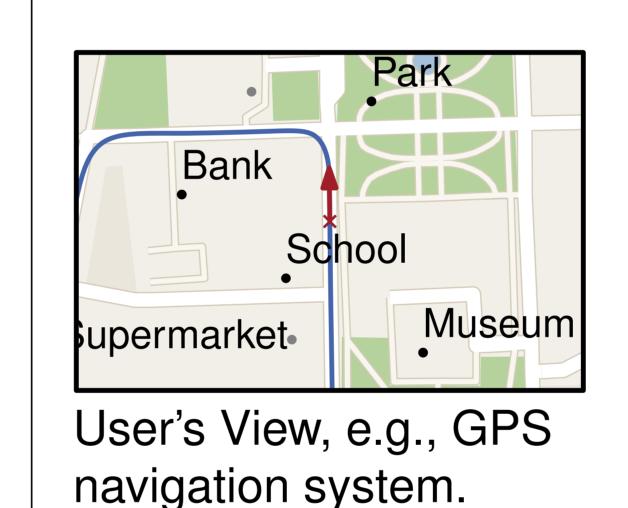
Boundary Labeling



Two-Sided Boundary Labeling with **Adjacent Sides:**

Rectangular section R of map. Labels in form of uniform rectangles

Trajectory-Based Map Labeling





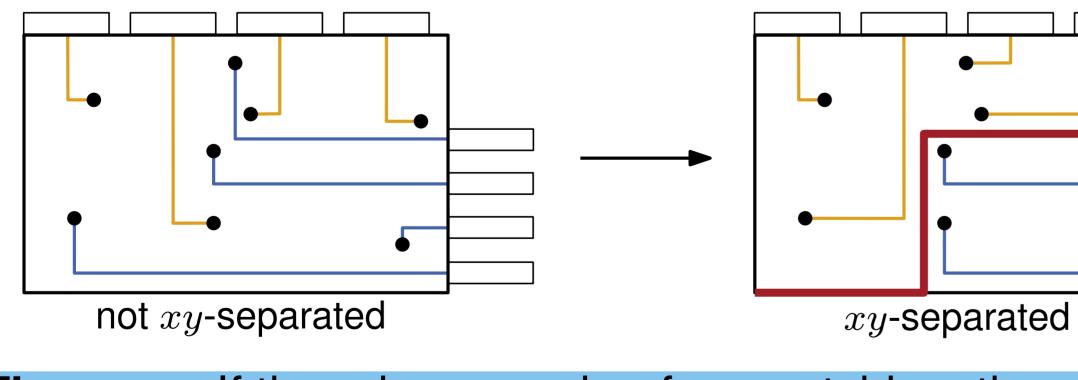
along two adjacent sides of R.

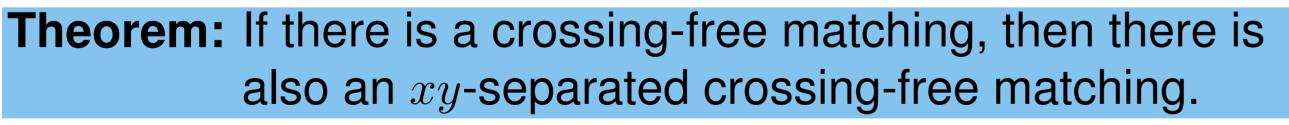
Sites on R in form of points.

Goal: Find optimal matching between labels and sites connecting them with *leaders*.

- **Requirements:** Minimize badness function, e.g., total length of leaders.
 - Leaders may not intersect each other, i.e., find crossing-free matching.

Structure of Solution:





Approach:



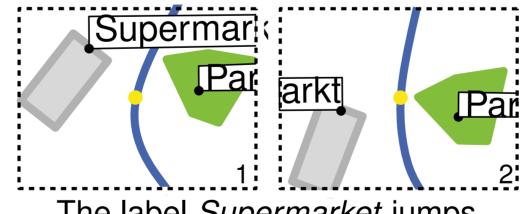
General View

Goal: Compute an optimal consistent labeling of a map along a given smooth trajectory.

Requirements: Maximize over all views the number of displayed labels. Superniarket Displayed labels may not overlap. Satisfy consistency criteria.

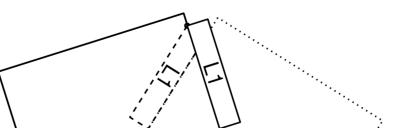
Consistency Criteria:

No 'jumping' labels.

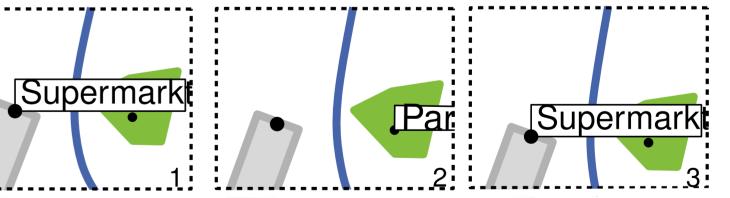


The label Supermarket jumps in the second frame.

Approach:



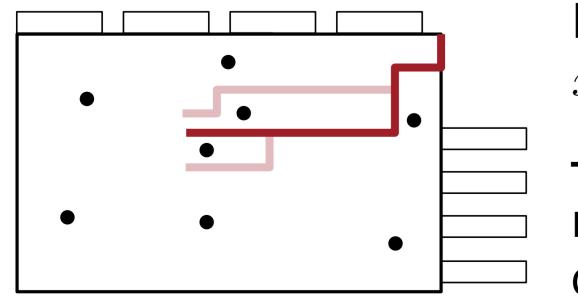
No 'flickering' labels.



The labels *Supermarket* and *Park* are displayed alternately.

Abstraction of Problem

Represent intersection between viewport and label by time intervals.



Dynamic program finds xy-separating curve in $O(n^2)$ time.

Crossing-free matching minimizing the total leader length can be computed in $O(n^2)$ time.

Joint Work with: Philipp Kindermann, Ignaz Rutter, Marcus Schaefer, André Schulz, and Alexander Wolff.

Publication: Two-Sided Boundary Labeling with Adjacent Sides. In: Algorithms and Data Structures, 13th International Symposium (WADS'13), Lecture Notes in Computer Science. Springer, 2013.

Time: L2 ⊢ L3 J Label is contained in viewport. Labels overlap.

Represent intersection between two labels by time intervals. Developed algorithms require only time intervals. General Case: NP-complete, approximation algorithms. **Display** $\leq \mathbf{k}$ labels at the same time: (Reduction of displayed information.) Polynomial solvable with dyn. program.

Joint Work with: Andreas Gemsa and Martin Nöllenburg.

KIT – University of the State of Baden-Wuerttemberg and National Laboratory of the Helmholtz Association

www.kit.edu